

## Re: "Theorem" in Mendelson ?

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- *From:* "MoeBlee" <jazzmobe@xxxxxxxxxxx>
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Question wrote:

Also, Mendelson makes the non-necessity that proofs be derived from axioms on the same page I quoted ("Introduction of Mathematical Logic," 4th ed. p. 34-5) where he reiterates, but even more clearly:

"A wf  $C$  is said to be a consequence in  $L$  of a set  $T$  [no better symbol for  $\gamma$ ] of wfs if and only if there is a sequence  $B_1, \dots, B_k$  of wfs such that  $C$  is  $B_k$  and, for each  $i$ , either  $B_i$  is an axiom or  $B_i$  is in  $T$ , or  $B_i$  is a direct consequence of some rule of inference of some of the preceding wfs in the sequence. Such a sequence is called a proof (or deduction) of  $C$  from  $T$ ."

Yes! That is exactly what I said I bet Mendelson provides. That is DIFFERENT from 'proof in  $L$ '.

He is given two definitions, each for a different relation. Or we can even do it four ways:

(1) A proof of in  $L$

(1a) A proof of  $C$  in  $L$

(2) A proof in  $L$  from  $T$

(2a) A proof of  $C$  in  $L$  from  $T$

And (1) is just 'a proof in  $L$  from the empty set'

And (1a) is just 'a proof of  $C$  in  $L$  from the empty set'

That is all very standard treatment and Mendelson gives it just fine. You can't just gloss over the difference when he talks about a proof in  $L$  and also talks about a proof in  $L$  from  $T$ .

There is no contradiction or problem with Mendelson. You just have to keep in mind the difference: proof in  $L$  vs. proof in  $L$  from  $T$ .

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And my earlier meta-proof that only tautologies may appear on lines in 'a proof in L' is correct. And your rebuttal was incorrect, since, for example, your proposed counterexample had  $B_1$  that is not an axiom and does not come from previous lines of the proof by a rule of inference.

On the other hand, your example IS a proof in L from T, where T is the set of premises, i.e., in this instance,  $\{B_1\}$ . That is, your non-tautologous  $B_1$  is okay as a line in a proof of L from T, since  $B_1$  is a member of T. But your non-tautologous  $B_1$  is not okay as a line in a proof from L, since in just a proof from L there IS NO T from which to draw premises and put them on lines in a proof.

A proof in L is something ensured to generate just tautologies.

A proof in L from T is something ensured to generate valid arguments – conclusions (that might not be tautologies) that are tautologically implied by premises (that might not themselves be tautologies).

And a proof in L is just a proof in L from the empty set.

Anyway, I see from your next paragraph (not quoted here) that you do understand something of this difference. But I want to leave my remarks here since I think they might help to encapsulate this crucial difference.

MoeBlee

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