

Re: Compact subsets of $\{0,1\}^{\mathbb{N}}$

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Summary of proof and comments presented by D.C. Ullrich regards

closed without isolated points, nonnull K subset $S = \{0,1\}^{\mathbb{N}}$
 $\implies K$ homeomorphic S

> If a is a finite sequence of 0's and 1's let S_a be the set of
> all elements of S that begin with a . Note that we're including
> the case where a is the empty sequence, ie a sequence of length 0:
> if a is the empty sequence then $S_a = S$.

> Let B be the tree of all finite sequences of 0's and 1's, and let
> T be the tree of all finite sequences a such that $S_a \cap K$
> is nonempty.

> Define a map F from B onto T by induction:

> If a is the empty sequence let $F(a) = a$.
> For each a in B , of length n , let a_0 and a_1 be the two extensions
> of a of length $n + 1$ (a_j consists of a followed by a j .)
> Now if a is in B and $F(a)$ has already been defined, define
> $F(a_0)$ and $F(a_1)$ as follows: Let $b = F(a)$.

Since K has no isolated points and S_b is open, $K \cap S_b$ must contain at least two points. So b has at least two extensions in T . Hence there exists an extension b' of b , with $\text{length}(b') = m$, say, such that

- (i) b' is the only extension of b in T of length m and
- (ii) b'_0 and b'_1 are both elements of T .

Define $F(a_0) = b'_0$ and $F(a_1) = b'_1$.

(Less formally: Since $K \cap S_b$ contains more than one point, the node b in T must eventually split in two, at some level below the level of b . Map a_0 and a_1 to the two nodes in T at the first level where b splits.)

> It follows that if a' is an extension of a then $F(a')$ is an
> extension of $F(a)$, so F induces a map $f: S \rightarrow S$ in a natural

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> way (given a in S , $f(a)$ is the unique element of S such that
> f (every finite initial segment of a) is an initial segment of $f(a)$).

> It's easy to see that f is continuous (given a and n there exists
> m such that if a' agrees with a in the first m places then $f(a')$
> agrees with $f(a)$ in the first n places).

f is 1-1: Say b, c are in S and $b \neq c$. There is a first place at which b and c differ, which says that there is a (possibly zero-length) finite string d such that d_0 is an initial segment of b and d_1 is an initial segment of c , or vice versa. Now by the construction of F , there exists e such that $F(d_0) = e_0$ and $F(d_1) = e_1$, so $F(d_0) \neq F(d_1)$. But $F(d_0)$ is an initial segment of $f(b)$ and $F(d_1)$ is an initial segment of $f(c)$, so $f(b) \neq f(c)$.

f maps S into K : Say a is in S . Let's say $a[:n]$ is the finite string consisting of the first n elements of a . Now by construction for every n there exists b^n in T such that $F(a[:n]) = b^n$, note that the length of b^n is at least n . Since b^n is in T , for every n there exists k^n in K such that b^n is an initial segment of k^n .

Now by construction $f(a)$ agrees with b^n in the first n places, and b^n also agrees with k^n in the first n places, so $f(a)$ agrees with k^n in the first n places. This says that k^n converges to $f(a)$, and hence $f(a)$ is in K since K is closed.

f maps S onto K : Let's say that a finite string t in T is an n -sibling if t has length n and there exists t' in T such that $t' \neq t$, t' has length n , and t and t' agree in the first $n-1$ places. Say that t in T is a sibling if it is the empty string or an n -sibling for some n . By construction the range of F is exactly the set of all siblings.

Now suppose that k is in K . Since K has no isolated points there exists a cofinal sequence of initial segments of k , each of which is a sibling. Since each sibling is in the range of F it follows that k is in the range of f .

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The fact that f maps S onto K is exactly the part that I originally said I hoped was obvious, because a formal explanation might be a little long. Ok, if it's obvious to me then I should be able to write down a formal explanation – here we go (I'll give an illustration of why every element of K starts with $F(0)$ or $F(1)$ after the proof):

Borrowing notation from Python, if $n \leq \text{length}(t)$ let's say

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that $t[:n]$ is the initial segment of t of length n . Now say t' is a child of t if $n = \text{length}(t) = \text{length}(t') - 1$ and $t'[:n] = t$. This "child of" relation makes T into a tree in which each node has one or two children.

Say t' is a descendant of t if t' is a child of t or a child of a child of t or etc. Say t' starts with t if $n = \text{length}(t) \leq \text{length}(t')$ and $t'[:n] = t$. So of course t' starts with t if and only if t' is a descendant of t ; in the one case we're thinking about actual sequences and in the other we're thinking about the tree structure.

Ok. The first fact is that every element of T has more than one descendant. This is where the fact that K has no isolated points is used: Say t is an element of T , choose k in K so that k starts with t (the existence of such a t is exactly the definition of T), and consider a sequence of distinct points of K that converge to k . By the definition of convergence, all but finitely many of those points of K start with t ; in particular there exist at least two elements of K that start with t , which says that t has at least two descendants.

So every element of T has at least two descendants, but of course an element of T can have only one child. Say t is in T . Define $\text{split}(t) = t'$, where t' is the element of T of minimal length such that t' is a descendant of t and t' has two children.

As before, let a_0 and a_1 be the two sequences consisting of a followed by a 0 or 1, respectively. The definition of F was this: $F(\text{empty sequence}) = \text{empty sequence}$, and $F(a_0) = \text{split}(F(a))_0$, $F(a_1) = \text{split}(F(a))_1$.

It follows that the range of F is exactly the empty sequence plus all siblings in T : one direction is immediate, and for the other direction, given t in T which is a sibling, trace the chain of ancestors up the tree until you come to an ancestor of shorter length which is also a sibling or the empty sequence. (If t^\wedge is a proper ancestor of t of maximal length such that t^\wedge is also a sibling then t^\wedge is in the range of F by induction on the length and then the definition of F shows that t is in the range of F . On the other hand, if no proper ancestor of t is a sibling then $t = \text{split}(\text{empty sequence})$ and it also follows that t is in the range of F .)

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Now fix k in K . The second key fact is this:
infinitely many initial segments of k are siblings.
(We need to show that for every N there exists $n \geq N$ such that $k[:n]$ is a sibling. Fix n . Since k is not an isolated point of K there exists k' in K with $k' \triangleleft k$ such that $k'[:N] = k[:N]$. Since $k' \triangleleft k$ there exists $n > N$ such that $k'[n] \triangleleft k[n]$ (where $s[n]$ denotes the n -th term of s . Choose n minimal with this property; then $k[:n]$ is a sibling, because $k[:n]$ and $k'[:n]$ are unequal but they are both children of $k[:n-1]$.)

So every element of K has infinitely many initial segments which are siblings. Hence every k in K has infinitely many initial segments which are in the range of F , and hence k is in the range of f .

Why every element of K starts with $F(0)$ or $F(1)$:
It could happen, say, that every element of K starts with 01101, but not every element of K starts with 011010 and not every element of K starts with 011011. If that happens then every element of K starts with 011010 _or_ 011011, and the definitions show that 01101 = split(empty sequence), hence $F(0) = 011010$ and $F(1) = 011011$.

Another way of looking at the whole thing – much less formal, but it may make you realize suddenly why it's all obvious:
Draw a picture of the tree T : An infinite tree where every node has one or two children and every node has more than one descendant. Any time you see a chain of nodes each with only one child, imagine collapsing that chain into a single node in a new tree. Then that new tree is a complete infinite binary tree – any two complete infinite binary trees are obviously isomorphic, and the F above (translated to our new tree) is just the obvious isomorphism.

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