

# Re: Computability

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- *From:* quasi <[quasi@xxxxxxxx](mailto:quasi@xxxxxxxx)>
  - *Date:* Mon, 14 Jan 2008 20:39:50 -0500
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On Mon, 14 Jan 2008 16:42:38 -0800 (PST), The Dougster 22044 <[DGoncz@xxxxxxxxxxxxx](mailto:DGoncz@xxxxxxxxxxxxx)> wrote:

On Jan 14, 4:46 am, quasi <[qu...@xxxxxxxx](mailto:qu...@xxxxxxxx)> wrote:

On Mon, 14 Jan 2008 04:14:39 -0500, quasi <[qu...@xxxxxxxx](mailto:qu...@xxxxxxxx)> wrote:

On Sun, 13 Jan 2008 20:02:36 -0800 (PST), The Dougster 22044 <[DGo...@xxxxxxxxxxxxx](mailto:DGo...@xxxxxxxxxxxxx)> wrote:

If  $(x/y)^p \equiv -1$  then  $(x/y)^{2p} \equiv 1$  and  $o(x/y, z) = 2p$  But we lose information when we write  $o(x/y, z) = 2p$ . That doesn't specify that  $(x/y)^p \equiv -1$ .  $(x/y)^p$  could be congruent to  $x$  for all we know! It just doesn't say.

Right.

Let me try to restate your conjecture ...

For greater simplicity, I won't bother with the order function. Instead, I'll just use the divisibility relation.

Dougster's conjecture:

## Re: Computability

There do not exist positive integers  $x, y, z$  such that

(1)  $x < y < z < x+y$

(2)  $x, y, z$ , are pairwise coprime

(3)  $z - y$  is not a multiple of  $x$

(4) For some prime  $p > 2$ ,

$$\begin{aligned}x &| z^p - y^p \\y &| z^p - x^p \\z &| x^p + y^p\end{aligned}$$

Remarks:

(1) For  $p = 3$ , your conjecture holds for  $z \leq 1000$ .

(2) As you've previously noted, if a proof of your conjecture could be had, that would yield an instant proof of FLT, however the known truth of FLT does not appear to yield a proof of your conjecture.

(3) As far as trying to prove your conjecture, I would start with a fixed prime, for example  $p = 3$ .

quasi

Re: Computability

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In the future, for new posts, please start the subject line with  
a  
prefix of "— ". By using this convention, such posts can be  
quickly  
recognized as "non-spam". Thanks.

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Well, it appears that your conjecture fails.

Here's a counterexample:

$$(x,y,z) = (43, 638, 659)$$

with  $p = 7$ .

Well, I wrote a  $check(x,y,z)$  function just for this counterexample,  
and no, it doesn't check out. I am flagging five conditions:

PASS:  $x < y < z < x+y$

PASS:  $(x,y) = (y,z) = (z,x) = 1$

PASS:  $(o(x/y,z), o(z/x,y), o(z/y,x)) = (2p, p, p)$

PASS:  $(x/y)^p \bmod z == -1$

FAIL:  $x + y - z = 0$  (because  $x == y == z \bmod p$  here)

Well that's the \_new\_ condition.

That condition was added after I posted the counterexample.

The old conditions only specified the other tests.

Hm. Some days when I am riding my bicycle I seem to see an answer to  
this problem, a contradiction buried here that would prove FLT. Maybe  
one day it will come to me, or to one of you.

Not likely. Firstly, it's probably not true. On the other hand, if it  
is true, there's no reason to expect an elementary proof.

Bottom line — defer this quest until you know a \_lot\_ more math.  
Don't let it compete for time against the time needed to break through  
to more advanced levels.

Note,  $43 == 638 == 659 == 1 \bmod 7$ .

Right.

So the new condition, makes your conjecture harder to defeat.

However there may still be counterexamples, simply less of them.

If you keep on adding more conditions every time you are presented with counterexample, you can keep the conjecture alive forever, but it would be artificial.

Looking at it in terms of probabilities, assuming independence of all the conditions, the probability may now be very close to zero. Thus, for example, if the probability is  $1/2^{100}$ , you are unlikely to find such a counterexample, even though such an example almost surely exists.

I don't see any strong reason why your conjecture should hold.

Moreover, given that each of the previous versions has fallen to counterexamples, skepticism seems called for.

quasi

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A number of sci.math participants have agreed to start all new article using a subject line with a prefix of "— ". By using this convention, most such posts (but perhaps not all) can be quickly recognized as "non-spam". I encourage others to try the idea. It's just an experiment to see if we can create an effective way to quickly recognize a message as probably non-spam.

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